

A Parallel Implementation with a comparative study of GMRES(m) Preconditionned by Multiplicative Schwarz, Additive Schwartz and Schur Complement

Emmanuel Kamgnia* Guy Antoine Atenekeng Kahou† Bernard Philippe ‡
Masha Sosonkina§

The GMRES(m) method. In this talk, we consider the GMRES(m) method the large linear system

$$Ax = b \tag{1}$$

$A \in \mathbb{R}^{n \times n}$, $b \in \mathbb{R}^n$.

For the GMRES(m) to be efficient, system (1) needs be preconditioned. The computational kernel of each GMRES(m) iteration consists of the following two modules :

- $x \rightarrow Ax$: matrix vector product (total of m) ;
- construction of the orthogonal basis.

The two modules are followed by an inexpensive solution of linear least squares problem. We propose a specialized parallel implementation.

Representation of distributed linear system. In a distributed solution of $Ax = b$, two types of variables may be distinguished: (1) Interior variables coupled only with local points of sub-domain; (2) Inter-domain interface variables coupled with external as well as the local interface points. All the local vectors, the vector of unknowns x_i and right hand side b_i ($i = 1, \dots, p$), are then divided as follows: the sub-vector u_i and f_i of interior variables followed by the sub-vector y_i and g_i of inter-domain interface variables. The local matrix A_i residing in processor i is block-partitioned according to this splitting:

$$\begin{pmatrix} B_i & F_i \\ E_i & C_i \end{pmatrix}, \tag{2}$$

where B_i corresponds to interior points, F_i and E_i couple inter-domain interface points, and C_i corresponds to external points.

*Department of Computer Science, University of Yaound I, P.O Box 812 Yaound, Cameroon, e-mail: ekamgnia@uycdc.uninet.cm

†Department of Mathematics and Computer Science, University of Dschang, P.O. Box Dschang, Cameroon, e-mail: gaatenek@irisa.fr

‡INRIA/IRISA Rennes, e-mail: Bernard.Philippe@irisa.fr

§Ames Laboratory, Iowa State University, Ames, IA 50011 USA, e-mail: masha@scl.ameslab.gov.

Multiplicative Schwarz preconditioner in explicit formulation The external points of partition p_i , namely C_i , should be in only one partition p_{i+1} . In terms of graph partitioning, this means that two partitions (sub-domains) p_i and p_j are neighbors if and only if $|i - j| = 1$. It is proved [2] that, in this case, we can explicitly set matrix M^{-1} of splitting defined by multiplicative Schwarz as

$$M^{-1} = \bar{A}_p^{-1} \bar{C}_{p-1} \bar{A}_{p-1}^{-1} \bar{C}_{p-2} \cdots \bar{A}_2^{-1} \bar{C}_1 \bar{A}_1^{-1}, \quad (3)$$

where for $i = 1, \dots, p$ matrix \bar{A}_i corresponds to the block A_i and for $i = 1, \dots, p - 1$ and matrix \bar{C}_i to the overlap C_i when completed by identity in both cases. We can obtain parallelism in EFMS by pipelining the construction of the Krylov subspace. In our implementation of GMRES, this is done by decoupling the construction of Krylov subspace from the QR factorization in the Arnoldi process.

Additive Schwarz and distributed Schur complement preconditioning Additive Schwarz (AS) procedures are the simplest parallel preconditioners available. They are easily parallelized and incur communications only in the exchange of interface variables, which may be the same as communications during a distributed matrix vector product and must precede the update of the local residual vector. This vector is used as the right-hand side for the local system to find the local update. There are several options for solving the local system: a (sparse) direct solver, a standard preconditioned Krylov solver, or a forward-backward solution associated with an accurate ILU preconditioner [1].

Schur complement (SC) methods iterate on the inter-domain interface unknowns only, implicitly using interior unknowns as intermediate variables. SC systems are derived by eliminating the variables u_i from (2): Once the global SC system is (approximately) solved, *each* processor computes the u -part of the solution vector by solving the system $B_i u_i = f_i - F_i y_i$ obtained by substitution from (2) [5].

Experiment descriptions First, we will compare the accuracy that may be achieved by the methods under consideration. Thus, we strive to achieve the best accuracy possible for each method. In other words, we compare the final (unscaled) residual of the original system defined as $\|b - Ax\|$. Separately, we will present the convergence curves. The comparison of solution and preconditioner construction times is also to be undertaken. The memory requirements will be compared in terms of the number of non-zeros in the preconditioner. The number of matrix-vector multiplies will help analyze the efficiency of the methods.

References

- [1] Yousef Saad, Iterative methods for Sparse Linear Systems, Society for Industrial and Applied Mathematics, SIAM, 2003
- [2] K. Atenekeng and E. Kamgnia and B. Philippe, An explicit formulation of the multiplicative Schwarz preconditionner, INRIA, 2005
- [3] R. Sidje and B. Philippe, Parallel krylov subspace basis computation, *CARI*, 1994, pp.421-440.
- [4] Z. Bai and D. Hu and L. Reichel, A Newton basis GMRES implementation, *IMA Journal of Numerical Analysis*, 14, 1994.
- [5] Y. Saad and M. Sosonkina, Distributed Schur Complement techniques for general sparse linear systems. *SIAM J. Sci. Comput.*, 21(4):1337-1356, 1999.